

Gentzen's sequent calculus

1935.a. Gerhard Genzen defined 1st order formulas as *sequents*.

Sequent

$$A_1, \dots, A_m \vdash B_1, \dots, B_n$$

is equivalent to *1st order formula*

$$A_1 \wedge \dots \wedge A_m \rightarrow B_1 \vee \dots \vee B_n$$

where $m, n \geq 0$ and $A_1, \dots, A_m, B_1, \dots, B_n$ are formulas

Sequent lh formula A_1, \dots, A_m – *antecedent*,

rh formula B_1, \dots, B_n – *succedent*.

Antecedent: A_1, \dots, A_m represents formula $A_1 \wedge \dots \wedge A_m$

Succedent: B_1, \dots, B_n represents formula $B_1 \vee \dots \vee B_n$.

$m = 0$, means that antecedent formula is unconditionally true

$n = 0$, means empty disjunct and contradiction

Language

Sequents consist of formulas constructed using : $\neg, \wedge, \vee, \Rightarrow, \forall, \exists$

Axiom (scheme) $A \rightarrow A$

Derivation rules - elimination and structural rules.

Each rule formalizes some proof step

Notations

Upper case latin letters A, B, \dots denote formulas

x – bound variable

a – free variable

t – term

$\Gamma, \Phi, \Lambda, \Pi$ – conjunctive/disjunctive sequences of formulae

Inference rules I

Axiom

$$\frac{}{A \vdash A} \quad (I)$$

$$\frac{\Gamma \vdash \Delta, A \quad A, \Sigma \vdash \Pi}{\Gamma, \Sigma \vdash \Delta, \Pi} \quad (Cut)$$

$$\frac{\Gamma, A \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad (\wedge L_1)$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \vee B, \Delta} \quad (\vee R_1)$$

$$\frac{\Gamma, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad (\wedge L_2)$$

$$\frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \vee B, \Delta} \quad (\vee R_2)$$

Inference rules II

$$\frac{\Gamma, A \vdash \Delta \quad \Sigma, B \vdash \Pi}{\Gamma, \Sigma, A \vee B \vdash \Delta, \Pi} \quad (\vee L)$$

$$\frac{\Gamma \vdash A, \Delta \quad \Sigma \vdash B, \Pi}{\Gamma, \Sigma \vdash A \wedge B, \Delta, \Pi} \quad (\wedge R)$$

$$\frac{\Gamma \vdash A, \Delta \quad \Sigma, B \vdash \Pi}{\Gamma, \Sigma, A \rightarrow B \vdash \Delta, \Pi} \quad (\rightarrow L)$$

$$\frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \rightarrow B, \Delta} \quad (\rightarrow R)$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma, \neg A \vdash \Delta} \quad (\neg L)$$

$$\frac{\Gamma, A \vdash \Delta}{\Gamma \vdash \neg A, \Delta} \quad (\neg R)$$

Inference rules III

$$\frac{\Gamma, A[t/x] \vdash \Delta}{\Gamma, \forall x A \vdash \Delta} \quad (\forall L)$$

$$\frac{\Gamma \vdash A[y/x], \Delta}{\Gamma \vdash \forall x A, \Delta} \quad (\forall R)$$

$$\frac{\Gamma, A[y/x] \vdash \Delta}{\Gamma, \exists x A \vdash \Delta} \quad (\exists L)$$

$$\frac{\Gamma \vdash A[t/x], \Delta}{\Gamma \vdash \exists x A, \Delta} \quad (\exists R)$$

*In rules $\forall R$ and $\exists L$ the variable y must not occur free within Γ and Δ .
Alternatively, the variable y must not appear anywhere in the respective lower sequents.*

Structural rules

$$\frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} \quad (WL)$$

$$\frac{\Gamma \vdash \Delta}{\Gamma \vdash A, \Delta} \quad (WR)$$

$$\frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} \quad (CL)$$

$$\frac{\Gamma \vdash A, A, \Delta}{\Gamma \vdash A, \Delta} \quad (CR)$$

$$\frac{\Gamma_1, A, B, \Gamma_2 \vdash \Delta}{\Gamma_1, B, A, \Gamma_2 \vdash \Delta} \quad (PL)$$

$$\frac{\Gamma \vdash \Delta_1, A, B, \Delta_2}{\Gamma \vdash \Delta_1, B, A, \Delta_2} \quad (PR)$$

Inference Example

$$\begin{array}{c}
 \frac{}{B \vdash B} (I) \quad \frac{}{C \vdash C} (I) \\
 \hline
 \frac{B \vee C \vdash B, C}{B \vee C \vdash C \wedge B} (\vee L) \\
 \hline
 \frac{B \vee C \vdash C \wedge B}{B \vee C, \neg C \vdash B} (\wedge R) \\
 \hline
 \frac{B \vee C, \neg C \vdash B}{\neg A \vdash \neg A} (\neg L) \quad \frac{}{\neg A \vdash \neg A} (I) \\
 \hline
 \frac{(B \vee C), \neg C, (B \rightarrow \neg A) \vdash \neg A}{(B \vee C), \neg C, ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A} (\rightarrow L) \\
 \hline
 \frac{(B \vee C), \neg C, ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A}{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C), \neg C \vdash \neg A} (\wedge L_1) \\
 \hline
 \frac{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C), \neg C \vdash \neg A}{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A} (PL) \\
 \hline
 \frac{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A}{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A} (\wedge L_2) \\
 \hline
 \frac{A \vdash A}{\vdash \neg A, A} (I) \quad \frac{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A}{\vdash \neg A, A} (\neg R) \\
 \hline
 \frac{\vdash \neg A, A}{\vdash A, \neg A} (PR) \quad \frac{(B \vee C), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A}{\vdash A, \neg A} (CL) \\
 \hline
 \frac{\vdash A, \neg A}{((B \rightarrow \neg A) \wedge \neg C), (B \vee C) \vdash \neg A} (PR) \\
 \hline
 \frac{((B \rightarrow \neg A) \wedge \neg C), (B \vee C) \vdash \neg A}{((B \rightarrow \neg A) \wedge \neg C), (A \rightarrow (B \vee C)) \vdash \neg A, \neg A} (\rightarrow L) \\
 \hline
 \frac{((B \rightarrow \neg A) \wedge \neg C), (A \rightarrow (B \vee C)) \vdash \neg A, \neg A}{((B \rightarrow \neg A) \wedge \neg C), (A \rightarrow (B \vee C)) \vdash \neg A} (CR) \\
 \hline
 \frac{((B \rightarrow \neg A) \wedge \neg C), (A \rightarrow (B \vee C)) \vdash \neg A}{(A \rightarrow (B \vee C)), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A} (PL) \\
 \hline
 \frac{(A \rightarrow (B \vee C)), ((B \rightarrow \neg A) \wedge \neg C) \vdash \neg A}{(A \rightarrow (B \vee C)) \vdash (((B \rightarrow \neg A) \wedge \neg C) \rightarrow \neg A)} (\rightarrow R) \\
 \hline
 \frac{(A \rightarrow (B \vee C)) \vdash (((B \rightarrow \neg A) \wedge \neg C) \rightarrow \neg A)}{\vdash ((A \rightarrow (B \vee C)) \rightarrow (((B \rightarrow \neg A) \wedge \neg C) \rightarrow \neg A))} (\rightarrow R)
 \end{array}$$